#### Five-Colour Theorem and Beyond

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Four-Colour Theorem.

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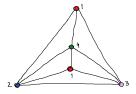
# Four-Colour Theorem and its controversy

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#### Four-Colour Theorem\*

Every planar graph can be properly coloured with four colours.



K. Appel, W. Haken (1977), Every Planar Map is Four Colorable Part I. Discharging, III. J. Math. 21: 429–490. K. Appel, W. Haken, J. Koch (1977), Every Planar Map is Four Colorable Part II. Reducibility, III. J. Math. 21: 491–567.

[2] K. Appel, W. Haken (1989), Every Planar Map is Four-Colorable, AMS.

[3] N. Robertson, D.P. Sanders, P. Seymour, R. Thomas (1997), The Four-Colour Theorem, JCTB 70: 2-44.

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5-C-T

# List colouring

List of admissible colours:  $\forall v \in V(G)$ :  $L(v) \subset \{1, 2, 3, ...\}$ List colouring  $\phi$ :

- $\forall v \in V(G): \phi(v) \in L(v)$
- $\forall uv \in E(G): \phi(u) \neq \phi(v)$

*k*-list-colouring: *L*-colouring where each vertex has k admissible colours.



Not every bipartite graph is always 2-list-colourable (i.e. not 2-choosable).

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5-C-T

### **5-Colour Theorem**



5-Colour Theorem (Thomassen, 1994). Planar graphs are 5-list-colourable.

Remark: Not all planar graphs are 4-choosable (Voigt, 1993).

# **Book proof**

Theorem. G near-triangulation

- ▶  $|L(a)|, |L(b)| \ge 1$  for two adjacent vertices on the infinite face
- $|L(u)| \ge 3$  for other vertices on the infinite face
- $|L(u)| \ge 5$  for the vertices not on the infinite face.

Then G is L-colourable.

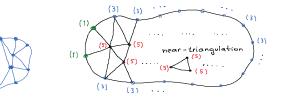


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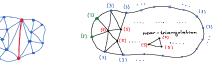
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A near-triangulation

Number of admissible colors

### Case 1 - Chords



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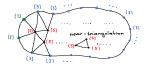


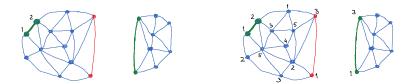




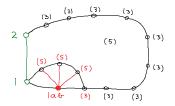
## Case 1 - Chords



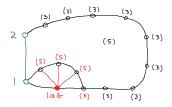


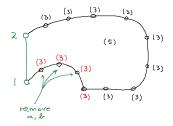


## Case 2 - No chords

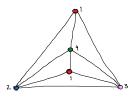


## Case 2 - No chords





# **Precolouring extension**

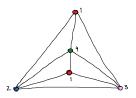


Two red vertices are coloured the same in every 4-colouring.

#### Question (Thomassen, 1990's)

► Precoloured vertices X of a planar graph, dist(x, y) ≥ 100, ∀x, y ∈ X. Can we extend to a 5-colouring?

# **Precolouring extension**



Two red vertices are coloured the same in every 4-colouring.

#### Question (Thomassen, 1990's)

- ▶ Precoloured vertices X of a planar graph,  $dist(x, y) \ge 100$ ,  $\forall x, y \in X$ . Can we extend to a 5-colouring?
- ▶ Precoloured vertices X of a planar graph,  $dist(x, y) \ge 10^{10}$ ,  $\forall x, y \in X$ . Can we extend to a 5-colouring?

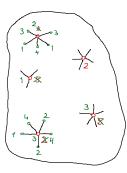
#### Not for 4-colouring!!

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#### Theorem (Albertson, 1998).

*G* 4-colourable,  $X \subset V(G)$  at distance  $\geq$  4 from each other. Then every 5-colouring of *X* extends to a 5-colouring of the whole graph.



# **Albertson's Conjecture**



Curious about the list-coloring version, Mike Albertson asked a question that became known as *Albertson's Conjecture*.

#### Question (Albertson, 1998). *G* planar, $X \subset V(G)$ at distance $\geq 10^{10}$ from each other, $|L(v)| \geq 5$ , $\forall v \in V(G)$ . Is it true that every *L*-colouring of *X* extends to an *L*-colouring of the whole graph.

### Some other extensions

There are other relaxations where 5-coloring results exist:

- Graphs drawn with crossings (far apart from each other).
- Precolored vertices (far apart).
- Locally planar graphs (arbitrary surfaces), no short non-contractible cycles.
- Longer precolored path on the infinite face (3 vertices).
- Precolored edges or triangles . . .

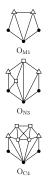
### **Recent results for list-colourings**

- Locally planar graphs (arbitrary surfaces), no short non-contractible cycles (DeVos, Kawarabayashi, M., 2008).
- Crossings at distance  $\geq$  19 from each other (Dvorak, Lidicky, M.)
- Albertson's Conjecture (Dvorak, Lidicky, M., Postle).
- Any combination of the above ingredients.

## **Crossings far apart**

- Precoloured path  $P \subset H$  on the outer face with up to 4 vertices
- Special subgraphs: Crossings, vertices with only 4 colours, 3-3 edges, far apart dist(A, B) ≥ r(A) + r(B) + 7
- Every obstruction is colourable
- 1/3/4/5 available colours.

# Obstructions



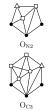








 $O_{\rm N1}$ 











 $O_{\rm C5}$ 





## Some more ingredients

*T*-critical w.r.t. *L*:  $T \subset G$ , s.t.  $\forall e \in E(G) \setminus E(T) : \exists L$ -colouring of *T* that extends to G - e but does not extend to *G*.

Lemma. If G is T-critical (w.r.t. L), then

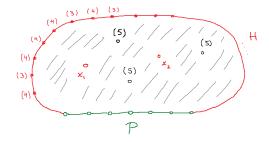
$$\omega_{\mathcal{T},L}(G) \leq |\mathcal{T}| - \frac{|\mathcal{V}(G) \setminus \mathcal{V}(\mathcal{T})|}{2|\mathcal{T}| + 2} - \frac{9}{2}.$$

where  $\omega_{T,L}(G)$  "counts" large faces and vertices in T with  $\geq 4$  colors.

Corollary. T connected subgraph of G, vertices in  $V(G) \setminus V(T)$  have  $\geq 5$  available colours (but no restrictions on T). If G is not L-colourable, then G contains a subgraph F on  $\leq 72|V(T)|^2$  vertices that is not L-colourable.

### **Basic setup**

- ► G plane graph, outer cycle H
- $P \subset H$  a path in H that is precoloured
- ▶ X (other) precoloured vertices, far apart
- ▶ 5 available colours for  $V(G) \setminus (V(H) \cup X)$
- ▶ 3/4/5 available colours for  $V(H) \setminus (V(P) \cup X)$ , no 3-3 edges!

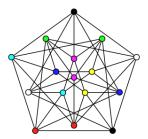


Theorem. (Extending a colouring of a path) If G is P-critical w.r.t. L, then  $\omega_{P,L}(G) \le |P| - 3$ .

This generalises Thomassen's theorem for |P| = 2 since  $\omega_{P,L}(G) \ge 0$ .

## **Outline of the proof of Albertson's Conjecture**

- ▶ Reduction to the case with |X| = 1: Using reduction techniques from our earlier paper (crossings far apart) and assuming we have a minimal counterexample.
- P ⊂ H a path in H that is precoloured. Since ω<sub>P,L</sub>(G) ≤ |P| − 3 and H has no 3-3 edges, we have a bound on the length of H (every vertex of H with list of size 4 contributes to ω).
- ▶ Reduction to the case when |X| = 1 and P has length ≤ 2. Part of the proof is computer supported.
- Here we characterise possible obstructions (infinite families) and then show that we can stay away from them.



#### Questions?